

# Truth and conditionals

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# Conditionals



# T-sentences

$$T(\ulcorner A \urcorner) \leftrightarrow A$$

# Reasoning

$A \therefore B$

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$$A, A \rightarrow B \vdash B$$

# Rejecting

$\neg \lambda$

# Weakening - I

$$\not\models T(\ulcorner \lambda \urcorner) \leftrightarrow \lambda$$

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$$\not\models T(\ulcorner \lambda \urcorner) \leftrightarrow \lambda$$

$$T(\ulcorner A \urcorner) \models A$$



## Weakening - II

$$A, A \rightarrow B \not\equiv B$$

# Weakening

$$\not\models T(\ulcorner \lambda \urcorner) \leftrightarrow \lambda$$

$$A, A \rightarrow B \not\models B$$



# Feferman objection



# Feferman objection

$$\begin{array}{|l} A \\ \hline \vdots \\ B \end{array} \quad A \supset B \quad \text{????}$$

# Conditionals

▷

# Conditionals

 $\supset$  $\supset \neq \rightarrow$

# Field



$$A \rightarrow B \models C \rightarrow A \rightarrow .C \rightarrow B$$

# Beall



$$A, A \rightarrow B \models B$$



# Several people

$$T(\ulcorner A \urcorner) \leftrightarrow A$$

## Field again, sort of

$$DA =_{Df} A \& \sim (A \rightarrow \sim A)$$

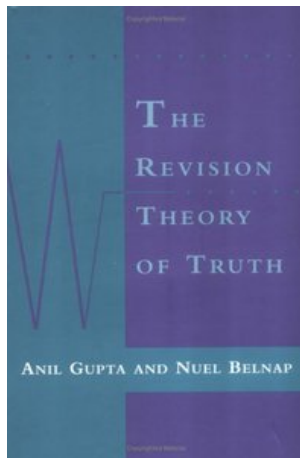
$$\dashv A = \sim D^* A$$

# Roles

Reasoning

Truth-theoretic features

# Revision theory



# Circular definitions

$$Gx =_{Df} A(x, G)$$

## Example

$$Gx =_{Df} (x = a \ \& \ \sim Gx) \vee (x = b \ \& \ Gx)$$

# Hypotheses

$$h \subseteq D$$

# Revising

$$GX =_{Df} A(x, G) \mapsto \delta$$



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$$GX =_{Df} A(x, G) \mapsto \delta$$

$$h, \delta(h), \delta(\delta(h)), \delta^3(h), \dots, \delta^\omega(h), \dots$$

# Example

$$Gx =_{Df} (x = a \ \& \ \sim Gx) \vee (x = b \ \& \ Gx)$$

	0	1	2	...
$\emptyset$	$\emptyset$	$\{a\}$	$\emptyset$	
$\{a\}$	$\{a\}$	$\emptyset$	$\{a\}$	
$\{b\}$	$\{b\}$	$\{a, b\}$	$\{b\}$	
$\{a, b\}$	$\{a, b\}$	$\{b\}$	$\{a, b\}$	

# Revision theory of truth

$$T(\ulcorner A_1 \urcorner) =_{Df} A_1$$

$$T(\ulcorner A_2 \urcorner) =_{Df} A_2$$

$$\vdots$$

$$T(\ulcorner A_n \urcorner) =_{Df} A_n$$

$$\vdots$$

# Classical logic and T-sentences

$$a = \ulcorner \sim Ta \urcorner$$

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$$\vdash_K \sim (A \equiv \sim A)$$

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$$a = \ulcorner \sim Ta \urcorner$$

$$\vdash_K \sim (A \equiv \sim A)$$

$$\not\vdash_{RT} Ta \equiv \sim Ta$$

# It gets worse

$$\vdash_{RT} \sim (Ta \equiv \sim Ta)$$

# Definitional equivalence

$$=_{Df} \neq \equiv$$



# Conditionals for revision theory

$$A \rightarrow B, A \leftarrow B$$

# Conditionals for revision theory

$$A \rightarrow B, A \leftarrow B$$

$$A \leftrightarrow B := (A \rightarrow B) \& (A \leftarrow B)$$

# New hypotheses

$$h \subseteq F \times V$$

# Semantics

$$M, v, h \models A \rightarrow B \Leftrightarrow M, v, h \not\models A \text{ or } \langle B, v \rangle \in_M h$$

$$M, v, h \models B \leftarrow A \Leftrightarrow \langle A, v \rangle \notin_M h \text{ or } M, v, h \models B$$

# Rules

$$\left| \begin{array}{l} A^{k+1} \\ \hline \vdots \\ B^k \end{array} \right| \quad A \rightarrow B^{k+1} \quad \rightarrow I \qquad \left| \begin{array}{l} A^{k+1} \\ A \rightarrow B^{k+1} \\ B^k \end{array} \right| \quad \rightarrow E$$

$$\left| \begin{array}{l} A^k \\ \hline \vdots \\ B^{k+1} \end{array} \right| \quad B \leftarrow A^{k+1} \quad \leftarrow I \qquad \left| \begin{array}{l} A^k \\ B \leftarrow A^{k+1} \\ B^{k+1} \end{array} \right| \quad \leftarrow E$$

# Features

$$\models_{RT+} T(\ulcorner A \urcorner) \leftrightarrow A$$

# Features

$$A =_{Df} B = A \leftrightarrow B$$

# Features

$$Gx =_{Df} A(x, C(Gx \leftrightarrow B))$$



# Logic - Sameness

- $(A \rightarrow C) \supset (A \& B \rightarrow C)$
- $(A \rightarrow B) \& (A \rightarrow C) \supset .A \rightarrow (B \& C)$
- $A \vee B \rightarrow C \supset .A \rightarrow C$
- $(A \rightarrow C) \& (B \rightarrow C) \supset .A \vee B \rightarrow C$
- $(\sim A \rightarrow B) \& (\sim A \rightarrow \sim B) \supset .A$

## Logic - Difference

- $\models ((C \leftarrow B) \leftarrow A) \equiv (C \leftarrow A \& B)$
- $\not\models (A \rightarrow (B \rightarrow C)) \supset A \& B \rightarrow C$
- $\not\models (A \& B \rightarrow C) \supset (A \rightarrow (B \rightarrow C))$
- $A \rightarrow (A \rightarrow B) \not\models A \rightarrow B$
- $(B \leftarrow A) \leftarrow A \models B \leftarrow A$

# Logic - Interaction

$$(A \rightarrow B) \equiv (\sim A \leftarrow \sim B)$$

# Flaws

# Flaws ...

# Flaws ...??

# Flaws ...??

$$\not\models A \rightarrow A$$

# Flaws ...??

$$\not\vdash A \rightarrow A$$

$$A \rightarrow B, B \rightarrow C \not\vdash A \rightarrow C$$



# Flaws ...??

$$\not\vdash A \rightarrow A$$



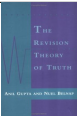
$$A \rightarrow B, B \rightarrow C \not\vdash A \rightarrow C$$

$$A \leftrightarrow B \not\vdash B \leftrightarrow A$$

# Seriously?

$\rightarrow$  ,  $\leftarrow$   $\neq$   $\Rightarrow$

# Roles revisited

	Reasoning	Truth
	$\rightarrow F$	$\rightarrow F$
	$\rightarrow BX$	$\rightarrow BX$
	$\supset$	$\rightarrow, \leftarrow$

# Too complicated

$$M, v, h \models A \rightarrow B \quad \Leftrightarrow \quad M, v, h \models A \text{ or } \langle B, v \rangle \in h$$

# Completeness

$$\models_{RT+}^{\mathcal{D}} A \Leftrightarrow \vdash_{RT+}^{\mathcal{D}} A$$

# Naturally fits into the revision theory

$$Gx =_{Df} A(x, G)$$

# Naturally fits into the revision theory

$$=_{Df} \neq \equiv$$

# Conclusions

- Distinguish roles conditionals play in our theories
- These roles can be used to motivate the addition of conditionals to logics
- Adding conditionals to revision theory fixes one of its problems
- These conditionals fill out the formal and philosophical picture of the revision theory
- Our earlier distinction can be used to defend these conditionals against objections



# Thank you

... to you, the audience.

... to Shunsuke Yatabe for inviting me.

... to James Shaw and the Pittsburgh philosophy department  
dissertation seminar for discussion.

... to Anil Gupta for the support, discussion, and many ideas and  
insights.

## Field quote

Field says the following of the strong Kleene material conditional.

*But while [the material conditional] does a passable job as a conditional in the presence of excluded middle, it is totally inadequate as a conditional without excluded middle: with  $\supset$  as one's candidate for  $\rightarrow$ , one wouldn't even get such elementary laws of the conditional as  $A \rightarrow A$ ,  $A \rightarrow (A \vee B)$ , or the inference from  $A \rightarrow B$  to  $(C \rightarrow A) \rightarrow (C \rightarrow B)$ . . . . The lack of a conditional (and also of a biconditional) cripples ordinary reasoning. Field (2008, 73)*

Field says that his conditional “enables us to come much closer to carrying out ordinary reasoning” than the strong Kleene material conditional does. Field (2008, 276)

## Feferman quote

“[N]othing like sustained ordinary reasoning can be carried on in [strong Kleene logic].” Feferman (1984, 95)

The whole quotation is emphasized in the original.

## Beall quote

“The question is whether we have detachable Tr-biconditionals (i.e., truth-biconditionals). If we do, then such biconditionals are not our usual material biconditionals, as noted above. I think that we do enjoy detachable Tr-biconditionals. . . .” (Beall, 2009, 26)

# Bibliography

Beall, J. (2009). *Spandrels of Truth*. Oxford University Press.

Feferman, S. (1984). Toward useful type-free theories. I. *Journal of Symbolic Logic*, 49(1):75–111.

Field, H. (2008). *Saving Truth from Paradox*. Oxford.